

Proofs!

Def'n A proof is a logically valid argument that establishes the truth of a proposition.

Proofs exist on a continuum from formal \rightarrow informal.

A formal proof:

- Uses only axioms (assumptions) and things previously proved
- Consists of a series of steps where each step is a valid logical inference from previous steps or assumptions.
- Is often expressed in formal notation.

An informal proof:

- Uses only axioms + things previously proved
- May omit or combine steps.
- Is often expressed in natural language.
- In principle could be expanded into a complete formal proof.
- Has other humans as its audience.

"Logical inference"

Rely on intuition!

eg. If $P \rightarrow Q$ is true and P is true, then Q is true.

(Why? intuitive; or make truth table; or show

$$((P \rightarrow Q) \wedge P) \rightarrow Q \equiv T$$

← practice for
learning goal
L3?

eg. If $A \wedge B$ is true, then A is true.

etc..

How To Prove Things

Just 3 easy steps!

- ① Translate the statement to be proved into propositional logic (using $\wedge, \vee, \neg, \rightarrow, \forall, \exists, \dots$)
- ② Write a proof outline corresponding to the propositional logic formula.
- ③ Use your intuition/insight/ingenuity/content knowledge to fill in the missing pieces.

$P \wedge Q$

To prove a conjunction $P \wedge Q$, prove P , and then prove Q .

We must show $P \wedge Q$, so we will prove both.

- | Proof of P
- | proof of Q

Therefore, $P \wedge Q$.

eg. Prove $3 < 5 \wedge 2 \times 8 = 16$.

Proof. We must show $3 < 5 \wedge 2 \times 8 = 16$, so we will prove both separately.

| First, $3 < 5$ because $3 + 2 = 5$.

| Second, $2 \times 8 = 16$ because $\begin{matrix} \circ & \circ & \circ & \circ & \circ & \circ \\ \circ & \circ & \circ & \circ & \circ & \circ \end{matrix}$.

Therefore, $3 < 5 \wedge 2 \times 8 = 16$.

$P \vee Q$

To prove $P \vee Q$, you have options:

- ① Prove P
- ② Prove Q
- ③ Use a proof by contradiction.

We must show $P \vee Q$, which we will do by proving P .

| Proof of P

So, since P is true, $P \vee Q$ must be true.

$P \rightarrow Q$

To prove $P \rightarrow Q$, suppose P is true, then prove Q in the hypothetical world where P is true.

We must show $P \rightarrow Q$, so suppose P .

| Proof of Q (using P)

Since Q is true under the supposition P , therefore $P \rightarrow Q$.

OR prove the contrapositive $\neg Q \rightarrow \neg P$.

We must show $P \rightarrow Q$, which we will do by proving the contrapositive, that is, $\neg Q \rightarrow \neg P$.

| Proof of $\neg Q \rightarrow \neg P$.

We proved $\neg Q \rightarrow \neg P$, which is equivalent to $P \rightarrow Q$.

$P \leftrightarrow Q$

To prove $P \leftrightarrow Q$, prove $(P \rightarrow Q) \wedge (Q \rightarrow P)$.

We must prove $P \leftrightarrow Q$, so we will prove both directions.

| (\rightarrow) proof of $P \rightarrow Q$.

| (\leftarrow) proof of $Q \rightarrow P$.

Therefore, since $P \rightarrow Q$ and $Q \rightarrow P$, $P \leftrightarrow Q$.

[7P] To prove $\neg P$:

① Use De Morgan laws to "push the \neg inwards" first.
eg. to prove $\neg(P \vee Q)$, prove $\neg P \wedge \neg Q$

② Prove $P \rightarrow F$. ($P \rightarrow F \equiv \neg P \vee F \equiv \neg P$).

Example- Write a proof outline for $P \rightarrow (Q \vee R)$.

To prove $P \rightarrow (Q \vee R)$, suppose P , then we will show $(Q \vee R)$.

We will show $Q \vee R$ by proving Q .

| Proof of Q (using P)

| Therefore, since Q is true, $Q \vee R$ is true.

Since we proved $Q \vee R$ under the supposition P , therefore $P \rightarrow (Q \vee R)$.

Example. Prove $(A \vee B) \rightarrow \neg C$ using a proof by contrapositive.

Proof. We will show $(A \vee B) \rightarrow \neg C$ via its contrapositive, that is,
 $C \rightarrow \neg(A \vee B)$. So suppose C ; we will show $\neg(A \vee B)$.

| $\neg(A \vee B) \equiv \neg A \wedge \neg B$, we will prove each separately.

| Part of $\neg A$ (using C)

| Proof of $\neg B$ (using C)

| Hence $\neg A \wedge \neg B$, that is, $\neg(A \vee B)$.

Therefore $C \rightarrow \neg(A \vee B)$, which is equivalent to $(A \vee B) \rightarrow \neg C$.

Recall $\text{Odd}(n) = \exists k \in \mathbb{Z}. n = 2k + 1.$

Example. Prove that if k is odd, so is k^2 .

$$(\forall k \in \mathbb{Z}). \text{Odd}(k) \rightarrow \text{Odd}(k^2)$$

Proof. To show the implication $\text{Odd}(k) \rightarrow \text{Odd}(k^2)$, suppose $\text{Odd}(k)$, that is, there is some integer j such that $k = 2j + 1$.
We must show $\text{Odd}(k^2)$, that is, we must find an integer p such that $k^2 = 2p + 1$.

KNOW

WANT.

To show $k^2 = 2p + 1$:

$$k^2 = (2j + 1)^2 = 4j^2 + 4j + 1 = 2(2j^2 + 2j) + 1.$$

So we can pick $p = 2j^2 + 2j$, which is an integer since j is.

Therefore, since we showed k^2 is odd under the assumption that k is odd, $\text{Odd}(k) \rightarrow \text{Odd}(k^2)$.

$\forall x \in D. P(x)$ To prove a forall:

① Prove $P(d)$ for an arbitrary d .

To prove $\forall x \in D. P(x)$, let d be an arbitrary element in domain D , we will prove $P(d)$.

| Proof of $P(d)$.

Since d was arbitrary, in fact $\forall x \in D. P(x)$.

② Use induction. (later).

$\exists x:D. P(x)$

To prove an exists statement:

① Pick a specific value in domain D , call it d , and prove $P(d)$. (eg. to prove $\exists n:\mathbb{N}. P(n)$, might prove $P(2)$).

② Use a proof by contradiction.

Ex. Prove $\forall n:\mathbb{Z}. \text{Odd}(3n+2) \rightarrow \text{Odd}(n)$.

Proof. (attempt)

Let z be an arbitrary integer; we will show $\text{Odd}(3z+2) \rightarrow \text{Odd}(z)$.

To show this implication, suppose $3z+2$ is odd; we will then show z is also odd.

$\text{Odd}(z) \equiv \exists k:\mathbb{Z}. z = 2k+1$. To prove this, we must give a specific value of k that makes $z = 2k+1$ true.

Since $3z+2$ is odd, there exists an integer j such that $3z+2 = 2j+1$. Solving for z ,

$$\Rightarrow 3z = 2j - 1$$

$$\Rightarrow z = \frac{2j-1}{3}$$

STUCK! Want $z = 2? + 1$ but this does not look like that.

Let's try contrapositive instead!

Ex. Prove $\forall n \in \mathbb{Z}. \text{Odd}(3n+2) \rightarrow \text{Odd}(n)$.

Proof.

Let n be an arbitrary integer.

To show $\text{Odd}(3n+2) \rightarrow \text{Odd}(n)$, we will use the contrapositive, namely, $\neg \text{Odd}(n) \rightarrow \neg \text{Odd}(3n+2)$, that is,

$\text{Even}(n) \rightarrow \text{Even}(3n+2)$. So

Suppose n is even. — KNOW

— WANT.

If n is even, $n = 2k$ for some integer k .

$$\begin{aligned} 3n+2 &= 3(2k) + 2 \\ &= 6k + 2 \\ &= 2(3k+1) \end{aligned}$$

Hence $3n+2$ is of the form $2 \times (\text{integer})$, so it is even.

We assume
 $\neg \text{odd} \equiv \text{Even}$
 $\neg \text{Even} \equiv \text{odd}$ —
we will prove later!