

Strong induction

Ex. Recall the Fibonacci sequence:

$$F_0 = 0$$

$$F_1 = 1$$

$$\underline{\underline{F_n = F_{n-1} + F_{n-2}}}$$

$$0, 1, 1, 2, 3, 5, 8, 13, \dots$$

$$1 \quad 2 \quad 4 \quad 8 \quad 16 \quad 32 \quad 64 \dots$$

Conjecture: $\forall n: \mathbb{N}. \underline{\underline{F_n \leq 2^{n-1}}}$. $\underbrace{\hspace{10em}}_{P(n)}$

Proof (attempted). Let's prove by induction. We must prove $P(0) \wedge (\forall k: \mathbb{N}. P(k) \rightarrow P(k+1))$.

• $P(0)$ says $F_0 \leq 2^{-1} \equiv 0 \leq \frac{1}{2} \checkmark$

• Now we must prove the induction step, $\forall k: \mathbb{N}. P(k) \rightarrow P(k+1)$.

So suppose k is an arbitrary natural # and suppose $P(k)$, that is, $F_k \leq 2^{k-1}$. We must show $P(k+1)$, that is, $F_{k+1} \leq 2^k$.

$$\begin{aligned} & F_{k+1} \\ &= F_k + F_{k-1} \\ &\leq 2^{k-1} + F_{k-1} \end{aligned}$$

{def'n of Fibonacci #'s}

X

$F_n = F_{n-1} + F_{n-2}$
only if $n \geq 2$.

Strong induction

Idea: when proving $P(k+1)$, we can assume $P(0), P(1), P(2), \dots, P(k)$ are all true, not just $P(k)$.

Assume $P(0), P(1), \dots, P(k)$ all true.

Formally:

$$\begin{aligned} P(0) \wedge (\forall k: \mathbb{N}. [\forall j: \mathbb{N}. j \leq k \rightarrow P(j)]) &\rightarrow P(k+1) \\ &\rightarrow (\forall n: \mathbb{N}. P(n)) \end{aligned}$$

Conjecture: $\forall n \in \mathbb{N}. F_n \leq 2^{n-1}$.

Proof: by strong induction.

- $P(0) \equiv F_0 \leq 2^{-1} \checkmark$
- $P(1) \equiv F_1 \leq 2^0 \equiv 1 \leq 1 \checkmark$
- Now let k be an arbitrary

natural number $k \geq 1$. We must

show $P(k+1)$ is true, if we

suppose $P(0), P(1), \dots, P(k)$ are true.

Need two

base cases

so the induction step can

"get off the ground"



$$\begin{aligned} & F_{k+1} \\ &= \quad \quad \quad \{ \text{defn of Fibonacci, since } k \geq 1. \} \\ & \leq \quad \quad \quad \{ P(k) \text{ and } P(k-1) \text{ are true} \}. \\ & \quad \quad \quad 2^{k-1} + 2^{k-2} \\ & < \quad \quad \quad \{ 2^{k-2} < 2^{k-1} \} \\ & \quad \quad \quad 2^{k-1} + 2^{k-1} \\ & = \quad \quad \quad \{ \text{algebra} \} \\ & \quad \quad \quad 2 \cdot 2^{k-1} \\ & = \quad \quad \quad \{ \text{algebra} \} \\ & \quad \quad \quad 2^k. \end{aligned}$$

Hence, $F_{k+1} \leq 2^k$.

Therefore, by strong induction, $F_n \leq 2^{n-1}$ for all $n \in \mathbb{N}$.

Note: strong + "weak" induction are actually logically equivalent.

Challenge: prove $\forall n \in \mathbb{N}. (3/2)^n < F_n$.

Ex: Fundamental Theorem of Arithmetic: every $n \geq 1$ can be written uniquely as a product of primes.

We will just prove every integer ≥ 2 can be written as a product of primes.

Proof. Let $P(n)$ be the proposition " n can be written as a product of primes." We will prove $\forall n \geq 2, P(n)$ by (strong?) induction.

- $P(2)$ says " 2 can be written as a prod. of primes": \checkmark
 2 is prime.
- Now let $k \geq 2$, and suppose $P(2), P(3), \dots, P(k)$ are all true. We must show $P(k+1)$.
 - $k+1$ is either prime or composite.
 - if it is prime, then $P(k+1)$ is true. since $k+1$ is a product of itself.
 - otherwise, there must be a, b such that $ab = k+1$ and $2 \leq a, b \leq k$. By our induction hypothesis, $P(a)$ and $P(b)$ are true, that is, a and b can be written as products of primes. Therefore so can $k+1 = ab$.

↪