

Logical Equivalences

Defn if P and Q always have the same truth value no matter what, we say they are logically equivalent and write $P \equiv Q$.

Core equivalences / algebraic rules for logic.

Principles:

1. Most equivalences have a "opposite world" version where we switch T/F and \wedge/\vee .
2. We can get intuition by thinking in terms of addition & multiplication:
 - T is kind of like 1
 - F is kind of like 0
 - \wedge is kind of like multiplication
 - \vee is kind of like addition (w/ max value of 2).

Name	equivalence	"opposite world" version.
Identity	$P \wedge T \equiv P$	$P \vee F \equiv P$
Annihilation	$P \wedge F \equiv F$	$P \vee T \equiv T$
Idempotence	$P \wedge P \equiv P$	$P \vee P \equiv P$
Commutativity	$P \wedge Q \equiv Q \wedge P$	$P \vee Q \equiv Q \vee P$
Associativity	$(P \wedge Q) \wedge R \equiv P \wedge (Q \wedge R)$	$(P \vee Q) \vee R \equiv P \vee (Q \vee R)$
Distributivity	$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)$	$P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)$
De Morgan	$\neg(P \wedge Q) \equiv \neg P \vee \neg Q$	$\neg(P \vee Q) \equiv \neg P \wedge \neg Q$
	Contradiction: $P \wedge \neg P \equiv F$	Law of Excluded Middle: $P \vee \neg P \equiv T$

Double negation elimination: $\neg(\neg P) \equiv P$

Implication $P \rightarrow Q \equiv \neg P \vee Q$

Iff $P \leftrightarrow Q \equiv (P \rightarrow Q) \wedge (Q \rightarrow P)$.

Example. Prove $(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$.

Strategy: use a series of equivalences

$$(p \rightarrow q) \wedge (p \rightarrow r) \equiv \dots \equiv \dots \equiv \dots \equiv p \rightarrow (q \wedge r).$$

relationship

$$\begin{aligned} & (p \rightarrow q) \wedge (p \rightarrow r) \\ \equiv & \quad \quad \quad \{ \text{Implication} \} \quad \quad \quad \} \leftarrow \text{reason why relationship holds} \\ & (\neg p \vee q) \wedge (\neg p \vee r) \\ \equiv & \quad \quad \quad \{ \text{Distributivity} \} \\ & \neg p \vee (q \wedge r) \\ \equiv & \quad \quad \quad \{ \text{Implication} \} \\ & p \rightarrow (q \wedge r) \end{aligned}$$

Example. Show $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$.

$$\begin{aligned} & \neg(p \vee (\neg p \wedge q)) \\ \equiv & \quad \quad \quad \{ \text{De Morgan} \} \\ & \neg p \wedge \neg(\neg p \wedge q) \\ \equiv & \quad \quad \quad \{ \text{De Morgan} \} \\ & \neg p \wedge (\neg(\neg p) \vee \neg q) \\ \equiv & \quad \quad \quad \{ 2 \times \text{negate} \} \\ & \neg p \wedge (p \vee \neg q) \\ \equiv & \quad \quad \quad \{ \text{Distributivity} \} \\ & (\neg p \wedge p) \vee (\neg p \wedge \neg q) \\ \equiv & \quad \quad \quad \{ \text{Contradiction} \} \\ & F \vee (\neg p \wedge \neg q) \\ \equiv & \quad \quad \quad \{ \text{Identity} \} \\ & \neg p \wedge \neg q \\ \equiv & \quad \quad \quad \{ \text{De Morgan} \} \\ & \neg(p \vee q) \end{aligned}$$