

Questions about graphs?

- # of paths $s \rightarrow t$? $\leftarrow$ later in semester
- Connected? $\leftarrow$ today
- # of edges?
- fastest/shortest/best path $s \rightarrow t$? $\leftarrow$ next week
- cycles? $\leftarrow$ HW
- tree? $\leftarrow$ today.

Q for today: is there a path $s \rightarrow t$? If so, what is it?

Idea: explore in one "direction" until getting stuck, then backtrack etc.

Mark where we've been to avoid getting caught in cycles.

Mark all vertices UNVISITED $\left\{ \begin{array}{l} \text{dictionary vertices} \rightarrow \text{bool} \\ \text{array of booleans} \\ \text{set} \end{array} \right\} \Theta(1) \text{ operations.}$

[parent $\leftarrow$ empty dict]

Depth-first search

DFS(G, u):

Mark u VISITED $\Theta(1)$

for each neighbor v of u :

if v is UNVISITED: $\Theta(1)$

[parent[v] $\leftarrow u$.] $\Theta(1)$

DFS(G, v)

assume we can list neighbors of u in $\Theta(\deg(u))$

Can answer:

- Is graph connected?
- Is s connected to t ?

To find a path $s \rightarrow t$, use the reversal of t , parent[t], parent[parent[t]], etc.

Time complexity?

- We look @ each vertex at most once. Once visited, we never look @ it again.
- We look @ every edge twice, once from each direction

Hence, overall time complexity is $\Theta(V + E)$

of vertices

of edges.