

Graph Q's?

- Path between 2 vertices? \leftarrow today
- Shortest path? \leftarrow Monday
- Acyclic? \rightarrow Tree? \leftarrow next weds.
- Connected? \leftarrow HW
- Number of edges?]
- Number of leaves?
- Number of cycles? \leftarrow !?

Depth-First Search (DFS)

Idea: explore as far as we can in one "direction", then back-track + try other directions.

Use "markers" to record where we've been, so we don't get stuck in a loop.

Mark all vertices UNVISITED
[parent \leftarrow empty dictionary]

$\left\{ \begin{array}{l} \text{list/array of booleans} \\ \text{dictionary vertices} \rightarrow \text{booleans} \\ \text{Set} \end{array} \right\}$ $\oplus(1)$ to mark and check degree of u .

DFS(G, u):

Mark u VISITED $\oplus(1)$

for each neighbor v of u : $\oplus(1)$
if v is UNVISITED: $\oplus(1)$

[parent[v] $\leftarrow u$.] $\oplus(1)$

DFS(G, v) \leftarrow

assume it takes $O(\deg(u))$ time to list neighbors of u .

we look at each vertex at most once.
 we look at each edge at most twice.

To check if s, t are connected: run $\text{DFS}(G, s)$, check if t was VISITED.

To find a path $s \rightarrow t$, run $\text{DFS}(G, s)$, then trace a path backwards from t to s : $t, \text{parent}[t], \text{parent}[\text{parent}[t]], \dots, s$.

Time complexity?

Time complexity of DFS is $O(V + E)$.

of vertices

of edges.