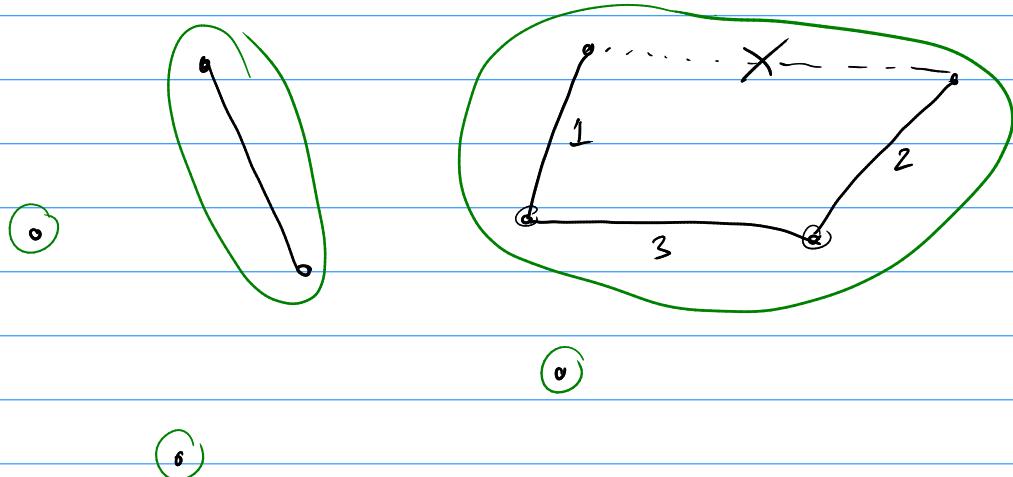


Kruskal's Algorithm

Sort the edges smallest \rightarrow biggest

For each edge:

Pick it if it doesn't make a cycle. \leftarrow



We want a data structure to keep track of sets of connected vertices. It should support the following operations:

- Union — union two sets containing 2 given vertices.
- Find — given a vertex, which set is it in?
- initialize — every vertex is in its own singleton set.

Often called a "disjoint set" data structure, aka "union-find".

Sort the edges smallest \rightarrow biggest $\leftarrow \Theta(E \lg V)$

For each edge (u, v) :

if $\text{find}(u) \neq \text{find}(v)$: $\leftarrow \Theta(E \cdot T_{\text{find}})$

pick edge (u, v)

$\text{union}(u, v)$.

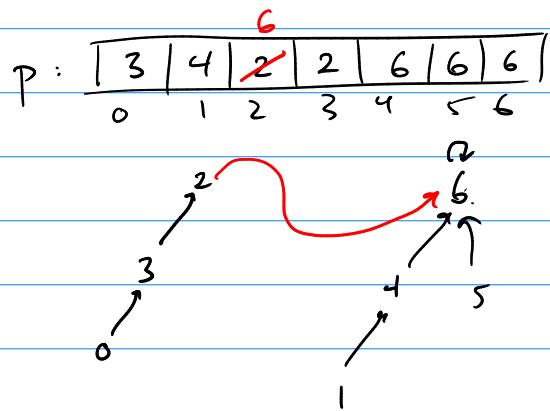
$\leftarrow \Theta(V \cdot T_{\text{union}})$

Overall: $\Theta(E \lg V + E \cdot T_{\text{find}} + V \cdot T_{\text{union}})$.

Goal: T_{find} and T_{union} $O(\lg V)$?

Implementing union-find.

Idea: keep a dictionary where every item points to another item which is in the same set. (its "parent").

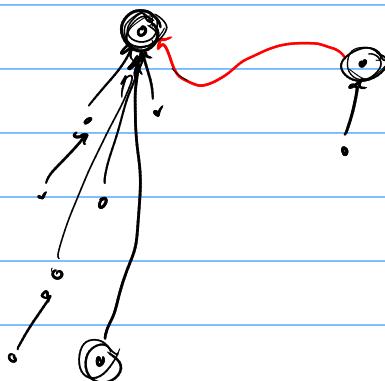


Each set is stored as a tree, with each vertex mapped to its parent.

The root of each tree is name of the set.

Union is now constant.

Find — $O(n)$ worst case? $O(\lg n)$ if we build trees carefully.



Idea: keep track of the height of each tree, by keeping a dict mapping nodes to height (height of non-root nodes doesn't matter).

When we union:

- Always make the node w/ smaller height child of the node w/ greater height.
- If the 2 trees have equal height, make either one the parent, increment its height.

Claim: using this scheme, find takes $O(\lg n)$.

- Path compression: when doing find, update everything along the path to point directly to the root.
 - Subsequent calls to find will be faster!

Fact: with path compression, find takes, on average, $\Theta(\lg^* n)$ time.

$\lg^* n = \# \text{ of times to repeat } \lg \text{ before reaching 1.}$

$$n = 2^{2^2} = 2^4 = 16, \quad \lg^*(n) = 3.$$

$$n = 2^{2^{2^2}} = 2^{2^4} = 2^{16} = 65,536, \quad \lg^*(n) = 4$$

$$n = 2^{2^{2^{2^2}}} = 2^{2^{2^4}} = 2^{16} = \text{HUGE!}$$

$\Rightarrow \lg^*(n)$ is ≤ 5 for all n s in our universe.